

Exploiting Variational Formulas for Quantum Relative Entropy

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joint work with Omar Fawzi and Marco Tomamichel

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Classical Relative Entropy

Definition (Relative Entropy)

For a positive measure Q on a finite set \mathcal{X} and a probability measure P on \mathcal{X} with $P \ll Q$, the relative entropy is defined as

$$D(P\|Q) := \sum_{x \in \mathcal{X}} P(x) \log \frac{P(x)}{Q(x)} \quad [\text{Kullback-Leibler 1951}], \quad (1)$$

where we understand $P(x) \log \frac{P(x)}{Q(x)} = 0$ whenever $P(x) = 0$.

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Definition (Rényi Relative Entropy)

For $\alpha \in (0, 1) \cup (1, \infty)$ the Rényi relative entropy is defined as

$$D_\alpha(P\|Q) := \frac{1}{\alpha - 1} \log \sum_{x \in \mathcal{X}} P(x)^\alpha Q(x)^{1-\alpha} \quad [\text{Rényi 1961}]. \quad (2)$$

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- What are the quantum extensions of $D(P\|Q)$ and $D_\alpha(P\|Q)$?

Measured Relative Entropy

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On a Hilbert space \mathcal{H} , for quantum states ρ and positive semi-definite operators σ the measured relative entropy is defined as

$$D^{\text{M}}(\rho||\sigma) := \sup_{(\mathcal{X}, M)} D(P_{\rho, M}||P_{\sigma, M}) \quad [\text{Donald 1986}], \quad (3)$$

where the optimization is over finite sets \mathcal{X} and positive operator valued measures (POVMs) M on \mathcal{X} , and $P_{\rho, M}(x) = \text{tr}[M(x)\rho]$.

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with the same premises as before.

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- Are there other quantum extensions?

Quantum Relative Entropy

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For the same setup as before, the quantum relative entropy is defined as

$$D(\rho||\sigma) := \text{tr}[\rho(\log \rho - \log \sigma)] \quad [\text{Umegaki 1962}] \quad (5)$$

if $\sigma \gg \rho$ and $+\infty$ if $\sigma \not\gg \rho$.

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Definition (Sandwiched Rényi Relative Entropy)

For $\alpha \in (0, 1) \cup (1, \infty)$, the Rényi relative entropy is defined as

$$D_\alpha(\rho\|\sigma) := \frac{1}{\alpha - 1} \log Q_\alpha(\rho\|\sigma) \quad \text{with} \quad Q_\alpha(\rho\|\sigma) := \text{tr} \left[\left(\sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right)^\alpha \right] \quad (6)$$

$$[\text{Müller-Lennert } et al. \text{ 2013, Wilde } et al. \text{ 2015}]. \quad (7)$$

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- How is $D^{\text{M}}(\rho\|\sigma)$ related to $D(\rho\|\sigma)$? How is $D_\alpha^{\text{M}}(\rho\|\sigma)$ related to $D_\alpha(\rho\|\sigma)$?

Variational Formulas for Relative Entropy I

- Variational Formula for Quantum Relative Entropy: [Petz 1988] we have

$$D(\rho||\sigma) = \sup_{\omega>0} \text{tr}[\rho \log \omega] + 1 - \text{tr}[\exp(\log \sigma + \log \omega)]. \quad (8)$$

Theorem (Variational Formula for Measured Relative Entropy)

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$$D^{\text{M}}(\rho\|\sigma) = \sup_{\omega>0} \text{tr}[\rho \log \omega] + 1 - \text{tr}[\sigma \omega]. \quad (9)$$

- By the Golden-Thompson inequality $\text{tr}[\log \sigma + \log \omega] \leq \text{tr}[\sigma \omega]$ we find:

Theorem (Achievability of Quantum Relative Entropy)

We have

$$D^{\text{M}}(\rho\|\sigma) \leq D(\rho\|\sigma) \quad \text{with equality if and only if } [\rho, \sigma] = 0. \quad (10)$$

Variational Formulas for Relative Entropy II

- Variational formula for sandwiched Rényi relative entropy: [Frank & Lieb 2013] for $Q_\alpha(\rho\|\sigma) = \exp((\alpha - 1)D_\alpha(\rho\|\sigma))$ we have

$$Q_\alpha(\rho\|\sigma) = \begin{cases} \inf_{\omega>0} \alpha \operatorname{tr}[\rho\omega] + (1 - \alpha) \operatorname{tr} \left[\left(\omega^{\frac{1}{2}} \sigma^{\frac{\alpha-1}{\alpha}} \omega^{\frac{1}{2}} \right)^{\frac{\alpha}{\alpha-1}} \right] & \text{for } \alpha \in (0, 1) \\ \sup_{\omega>0} \alpha \operatorname{tr}[\rho\omega] + (1 - \alpha) \operatorname{tr} \left[\left(\omega^{\frac{1}{2}} \sigma^{\frac{\alpha-1}{\alpha}} \omega^{\frac{1}{2}} \right)^{\frac{\alpha}{\alpha-1}} \right] & \text{for } \alpha \in (1, \infty). \end{cases} \quad (11)$$

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Theorem (Variational Formula for Measured Rényi Relative Entropy)

For $Q_\alpha^{\text{M}}(\rho\|\sigma) := \exp((\alpha - 1)D_\alpha^{\text{M}}(\rho\|\sigma))$ we have

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By the Araki-Lieb-Thirring inequality we have for $\alpha \in (1/2, \infty)$:

$$D_\alpha^{\text{M}}(\rho\|\sigma) \leq D_\alpha(\rho\|\sigma) \quad \text{with equality if and only if } [\rho, \sigma] = 0. \quad (13)$$

Application: Additivity in Quantum Information I

- We consider operational quantities of the form

$$\mathcal{M}(\rho) := \min_{\sigma \in \mathcal{C}} \mathbb{D}(\rho \| \sigma), \quad (14)$$

where $\mathbb{D}(\cdot \| \cdot)$ stands for any relative entropy, and \mathcal{C} denotes some convex, compact set of states.

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- separable states: relative entropy of entanglement [Vedral 1998]
- positive partial transpose states: Rains bound on entanglement distillation [Rains 2001]
- non-distillable states: bounds on entanglement distillation [Vedral 1999]
- quantum Markov states: robustness properties of these states [Linden *et al.* 2008]
- locally recoverable states: bounds on the conditional mutual information [Fawzi & Renner 2015]
- k -extendible states: bounds on squashed entanglement [Li & Winter 2014]

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- Question: What properties of the relative entropy translate to the measure \mathcal{M} ?
 - For example, all the relative entropies discussed are **super-additive** on tensor product states

$$\mathbb{D}(\rho_1 \otimes \rho_2 \| \sigma_1 \otimes \sigma_2) \geq \mathbb{D}(\rho_1 \| \sigma_1) + \mathbb{D}(\rho_2 \| \sigma_2). \quad (15)$$

Application: Additivity in Quantum Information II

- Super-additivity of \mathcal{M} on tensor product states:

$$\min_{\sigma_{12} \in \mathcal{C}_{12}} \mathbb{D}(\rho_1 \otimes \rho_2 \| \sigma_{12}) = \mathcal{M}(\rho_1 \otimes \rho_2) \stackrel{?}{\geq} \mathcal{M}(\rho_1) + \mathcal{M}(\rho_2) \quad (16)$$

$$= \min_{\sigma_1 \in \mathcal{C}_1} \mathbb{D}(\rho_1 \| \sigma_1) + \min_{\sigma_2 \in \mathcal{C}_2} \mathbb{D}(\rho_2 \| \sigma_2). \quad (17)$$

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- Idea: Use variational characterizations

$$\mathbb{D}(\rho \| \sigma) = \sup_{\omega > 0} f(\rho, \sigma, \omega) \text{ in order to write} \quad (18)$$

$$\mathcal{M}(\rho) = \min_{\sigma \in \mathcal{C}} \sup_{\omega > 0} f(\rho, \sigma, \omega) = \sup_{\omega > 0} \min_{\sigma \in \mathcal{C}} f(\rho, \sigma, \omega), \quad (19)$$

where we made use of Sion's minimax theorem for the last equality.

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where we made use of Sion's minimax theorem for the last equality.

- The minimization over $\sigma \in \mathcal{C}$ then often simplifies and is a convex or even semidefinite optimization. Using strong duality to rewrite this minimization as a maximization problem:

$$\min_{\sigma \in \mathcal{C}} f(\rho, \sigma, \omega) = \max_{\bar{\sigma} \in \bar{\mathcal{C}}} \bar{f}(\rho, \bar{\sigma}, \omega) \text{ leading to the expression} \quad (20)$$

$$\mathcal{M}(\rho) = \sup_{\omega > 0} \max_{\bar{\sigma} \in \bar{\mathcal{C}}} \bar{f}(\rho, \bar{\sigma}, \omega). \quad (21)$$

Application: Additivity in Quantum Information III

- Variational characterization (as on the last slide):

$$\mathcal{M}(\rho) = \sup_{\omega > 0} \max_{\bar{\sigma} \in \bar{\mathcal{C}}} \bar{f}(\rho, \bar{\sigma}, \omega). \quad (22)$$

Application: Additivity in Quantum Information III

- Variational characterization (as on the last slide):

$$\mathcal{M}(\rho) = \sup_{\omega > 0} \max_{\bar{\sigma} \in \bar{\mathcal{C}}} \bar{f}(\rho, \bar{\sigma}, \omega). \quad (22)$$

- The following two conditions on \bar{f} and $\bar{\mathcal{C}}$ imply **super-additivity** of \mathcal{M} :

- 1 \bar{f} is super-additive

$$\bar{f}(\rho_1 \otimes \rho_2, \bar{\sigma}_1 \otimes \bar{\sigma}_2, \omega_1 \otimes \omega_2) \geq \bar{f}(\rho_1, \bar{\sigma}_1, \omega_1) + \bar{f}(\rho_2, \bar{\sigma}_2, \omega_2) \quad (23)$$

- 2 the sets $\bar{\mathcal{C}}$ are closed under tensor products

$$\bar{\sigma}_1 \in \bar{\mathcal{C}}_1 \text{ and } \bar{\sigma}_2 \in \bar{\mathcal{C}}_2 \text{ imply that } \bar{\sigma}_1 \otimes \bar{\sigma}_2 \in \bar{\mathcal{C}}_{12} \quad (24)$$

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- Proof: For any $\omega_1, \omega_2 > 0$ and any $\bar{\sigma}_1 \in \bar{\mathcal{C}}_1, \bar{\sigma}_2 \in \bar{\mathcal{C}}_2$ we deduce that

$$\mathcal{M}(\rho_1 \otimes \rho_2) \geq \bar{f}(\rho_1 \otimes \rho_2, \bar{\sigma}_1 \otimes \bar{\sigma}_2, \omega_1 \otimes \omega_2) \geq \bar{f}(\rho_1, \bar{\sigma}_1, \omega_1) + \bar{f}(\rho_2, \bar{\sigma}_2, \omega_2) \quad (25)$$

Hence, the inequalities also hold true if we maximize over these variables, implying super-additivity. □

Example: Relative Entropy of Recovery I

- For any relative entropy $\mathbb{D}(\cdot \| \cdot)$ we are interested in the recovery quantity:

$$\mathbb{D}^{\text{rec}}(\rho_{AD} \| \sigma_{AE}) := \inf_{\Gamma_{E \rightarrow D}} \mathbb{D}(\rho_{AD} \| (\mathcal{I}_A \otimes \Gamma_{E \rightarrow D})(\sigma_{AE})). \quad (26)$$

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- Of interest because we have (the systems are understood as $D = BC$, $E = B$):

$$I(A : C|B) \geq \lim_{n \rightarrow \infty} \frac{1}{n} D^{\text{rec}}(\rho_{ABC}^{\otimes n} \| \rho_{AB}^{\otimes n}) \quad [\text{Brandao et al. 2015}]. \quad (27)$$

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Theorem (Super-Additivity of Measured Entropy of Recovery)

Let $\rho_{AD}, \tau_{A'D'}, \sigma_{AE}, \omega_{A'E'}$ be quantum states. Then, we have

$$D^{\text{M,rec}}(\rho_{AD} \otimes \tau_{A'D'} \| \sigma_{AE} \otimes \omega_{A'E'}) \geq D^{\text{M,rec}}(\rho_{AD} \| \sigma_{AE}) + D^{\text{M,rec}}(\tau_{A'D'} \| \omega_{A'E'}). \quad (28)$$

For $\alpha \in (0, 1) \cup (1, \infty)$, we also have

$$D_{\alpha}^{\text{M,rec}}(\rho_{AD} \otimes \tau_{A'D'} \| \sigma_{AE} \otimes \omega_{A'E'}) \geq D_{\alpha}^{\text{M,rec}}(\rho_{AD} \| \sigma_{AE}) + D_{\alpha}^{\text{M,rec}}(\tau_{A'D'} \| \omega_{A'E'}). \quad (29)$$

Example: Relative Entropy of Recovery II

- We conclude (without de Finetti reductions):

$$I(A : C|B) \geq \lim_{n \rightarrow \infty} \frac{1}{n} D^{\text{rec}} (\rho_{ABC}^{\otimes n} \| \rho_{AB}^{\otimes n}) \geq D^{\text{M,rec}} (\rho_{ABC} \| \rho_{AB}) . \quad (30)$$

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- The super-additivity is seen by the following dual variational characterization:

Lemma (Variational Representation of Measured Entropy of Recovery)

Let ρ_{AD}, σ_{AE} be quantum states, and let σ_{AEF} be a purification of σ_{AE} . Then, we have

$$\begin{aligned} D^{\text{M,rec}}(\rho_{AD} \| \sigma_{AE}) = \text{maximize:} & \quad \text{tr}[\rho_{AD} \log R_{AD}] \\ \text{subject to:} & \quad S_{AF} > 0, R_{AD} > 0 \\ & \quad 1_D \otimes S_{AF} \geq R_{AD} \otimes 1_F \\ & \quad \text{tr}[S_{AF} \sigma_{AF}] = 1. \end{aligned} \quad (31)$$

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- The characterization for the measured Rényi entropy of recovery $D_\alpha^{\text{M,rec}}$ is similar.
- The relative entropy of recovery D^{rec} and the Rényi entropy of recovery D_α^{rec} do not seem to be additive (but open).

Example: Relative Entropy of Recovery III

Lemma (Variational Representation of Entropy of Recovery)

For the same premises as before we have

$$\begin{aligned}
 D^{\text{rec}}(\rho_{AD} \parallel \sigma_{AE}) = \text{maximize:} & \quad \text{tr}[\rho_{AD} \log \rho_{AD}] - D^{\text{M}}(\rho_{AD} \parallel R_{AD}) \\
 \text{subject to:} & \quad S_{AF} > 0, R_{AD} > 0 \\
 & \quad \mathbf{1}_D \otimes S_{AF} \geq R_{AD} \otimes \mathbf{1}_F \\
 & \quad \text{tr}[S_{AF} \sigma_{AF}] = 1.
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Example: Relative Entropy of Recovery III

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For the same premises as before we have

$$\begin{aligned}
 D^{\text{rec}}(\rho_{AD} \parallel \sigma_{AE}) = \text{maximize:} & \quad \text{tr}[\rho_{AD} \log \rho_{AD}] - D^{\text{M}}(\rho_{AD} \parallel R_{AD}) \\
 \text{subject to:} & \quad S_{AF} > 0, R_{AD} > 0 \\
 & \quad 1_D \otimes S_{AF} \geq R_{AD} \otimes 1_F \\
 & \quad \text{tr}[S_{AF} \sigma_{AF}] = 1.
 \end{aligned} \tag{32}$$

- This is to be compared to the dual characterization of the measured entropy of recovery $D^{\text{M,rec}}(\rho_{AD} \parallel \sigma_{AE})$:

$$\begin{aligned}
 D^{\text{M,rec}}(\rho_{AD} \parallel \sigma_{AE}) = \text{maximize:} & \quad \text{tr}[\rho_{AD} \log \rho_{AD}] - D(\rho_{AD} \parallel R_{AD}) \\
 \text{subject to:} & \quad S_{AF} > 0, R_{AD} > 0 \\
 & \quad 1_D \otimes S_{AF} \geq R_{AD} \otimes 1_F \\
 & \quad \text{tr}[S_{AF} \sigma_{AF}] = 1.
 \end{aligned} \tag{33}$$

Conclusion

- Measured relative entropy is strictly smaller than quantum relative entropy.
- Variational formulas for quantum relative entropy.
- These formulas are useful tools for studying additivity problems in quantum information theory.
- Super-additivity of measured relative entropy of recovery.
- Additivity of relative entropy of recovery remains open (does not seem to be additive).