Quantum Computation (484)

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Indicator Function

We will search for a (single) element \mathbf{x}_0 which is identified via an indicator function $f:\{0,1\}^n \to \{0,1\}$:

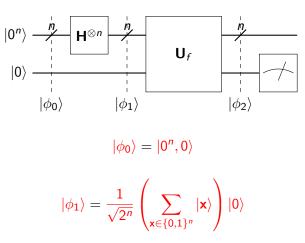
$$f(\mathbf{x}) = \begin{cases} 1 & \text{if } \mathbf{x} = \mathbf{x}_0 \\ 0 & \text{otherwise} \end{cases}$$

As usual, we represent this function as a unitary \mathbf{U}_f which maps $|\mathbf{x},y\rangle$ to $|\mathbf{x},f(\mathbf{x})\oplus y\rangle$. For example: to pick the element "10" out of $\{00,01,10,11\}$ we use

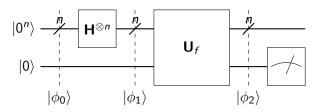
$$\mathbf{U}_f = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

Superposition

A first attempt could try to evaluate f or \mathbf{U}_f on all possible inputs.



Measuring the Final State

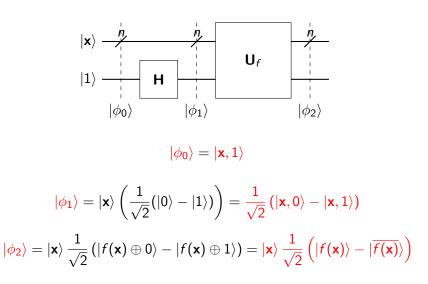


We can measure the top n qubits or the lower qubit of the final state:

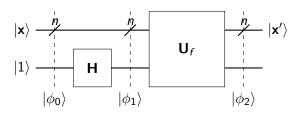
$$|\phi_2\rangle = \frac{1}{\sqrt{2^n}} \sum_{\mathbf{x} \in \{0,1\}^n} |\mathbf{x}, f(\mathbf{x})\rangle$$

If we measure the bottom qubit and if we are **lucky** we might measure $|1\rangle$. In this case the top qubits will collapse to the desired outcome \mathbf{x}_0 . Chances for this are just $\frac{1}{2^n}$.

Phase Inversion



Inverted State



$$|\phi_{2}\rangle = (-1)^{f(\mathbf{x})} |\mathbf{x}\rangle \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) = \begin{cases} -\frac{1}{\sqrt{2}} |\mathbf{x}\rangle (|0\rangle - |1\rangle) & \text{if } \mathbf{x} = \mathbf{x}_{0} \\ +\frac{1}{\sqrt{2}} |\mathbf{x}\rangle (|0\rangle - |1\rangle) & \text{if } \mathbf{x} \neq \mathbf{x}_{0} \end{cases}$$

Example: Searching for "10" in $\{00,01,10,11\}$. Starting with a uniform state $|\mathbf{x}\rangle = \frac{1}{2}(1,1,1,1)^T$ as top qubits, these become $|\mathbf{x}'\rangle = \frac{1}{2}(1,1,-1,1)^T$.

However, measuring $|\mathbf{x}'\rangle$ does not identify \mathbf{x}_0 .

Inversion about the Mean

Consider, for example, a sequence of numbers 53, 38, 17, 23, 79.

The average or mean of these is

$$a=42=\frac{53+38+17+23+79}{5}.$$

Replace every number v by its mirror image through a, i.e.

$$v' = a + (a - v)$$
$$= -v + 2a$$

In our example we get the new sequence 31, 46, 67, 61, 5 which has the same average a = 42.

Average Operator

Represent sequences as vectors \mathbf{v} , e.g. $\mathbf{v} = (53, 38, 17, 23, 79)^T$.

We can introduce an averaging operator for *n*-dimensional vectors

With this we can compute the average a of a sequence \mathbf{v} by:

$$\mathbf{Av} = (a, a, \dots, a)^T$$

In our example, $\mathbf{A}(53, 38, 17, 23, 79)^T = (42, 42, 42, 42, 42)^T$.

Mean Inversion Operator

We can also define a matrix representation of the inversion about the mean or average, which is in fact **unitary** (show!):

$$\mathbf{v}' = -\mathbf{v} + 2\mathbf{A}\mathbf{v}$$
$$= (-\mathbf{I} + 2\mathbf{A})\mathbf{v}$$

In our example the mean inversion operator is:

$$2\mathbf{A} - \mathbf{I} = \begin{pmatrix} \frac{2}{5} - 1 & \frac{2}{5} & \frac{2}{5} & \frac{2}{5} & \frac{2}{5} \\ \frac{2}{5} & \frac{2}{5} - 1 & \frac{2}{5} & \frac{2}{5} & \frac{2}{5} \\ \frac{2}{5} & \frac{2}{5} & \frac{2}{5} - 1 & \frac{2}{5} & \frac{2}{5} \\ \frac{2}{5} & \frac{2}{5} & \frac{2}{5} & \frac{2}{5} - 1 & \frac{2}{5} \\ \frac{2}{5} & \frac{2}{5} & \frac{2}{5} & \frac{2}{5} - 1 \end{pmatrix}$$

Thus we get: $(-\mathbf{I} + 2\mathbf{A})(53, 38, 17, 23, 79)^T = (31, 46, 67, 61, 5)^T$.

Amplitude Amplifying

Consider the following uniform vector of values:

$$\mathbf{v}_0 = (10, 10, 10, 10, 10)^T$$

Perform a phase inversion on the 4th element. The result is:

$$\mathbf{v}_1 = (10, 10, 10, -10, 10)^T$$

with average a = 6. Perform an **inversion about the mean** to get

$$\mathbf{v}_2 = (2, 2, 2, 22, 2)^T$$

followed by another phase inversion

$$\mathbf{v}_3 = (2, 2, 2, -22, 2)^T$$

and an **inversion about the mean** a = -2.8 to get:

$$\mathbf{v}_4 = (-7.6, -7.6, -7.6, 16.4, -7.6)^T$$

Grover's Algorithm

Use phase inversion and inversion about the mean to make it more likely to be **lucky**. These amplifications need to be applied $\sqrt{2^n}$ times for an n qubit register (i.e. 2^n search items).

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Step 1. Start with |0^n\rangle.
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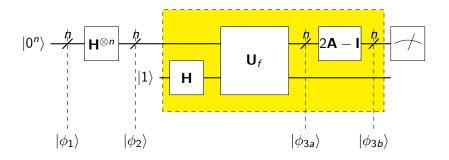
Step 2. Apply $\mathbf{H}^{\otimes n}$.

Step 3. Repeat $\sqrt{2^n}$ times:

Step 3a. Phase inversion $U_f(I \otimes H)$. Step 3b. Inversion about the mean -I + 2A.

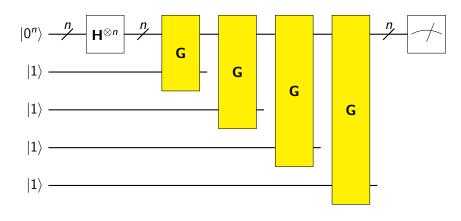
Step 4. Measurement.

Grover's Circuit

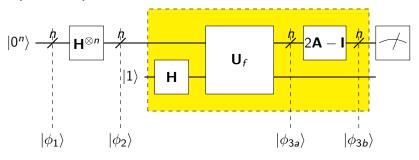


Classical search in an unordered array of size m needs m steps in the worst case and $\frac{m}{2}$ in the average. Grover's search algorithm will take \sqrt{m} steps.

Iterations of Grover's Gate



Example: Preparation



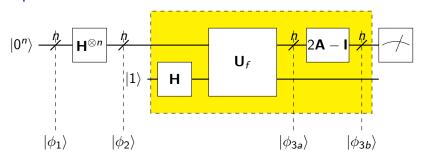
Find "110" in $\{000,001,010,011,100,101,110,111\}.$ We need $\sqrt{8}=2.8284$ iterations.

Preparation Step.

$$|\phi_1\rangle = (1,0,0,0,0,0,0,0)^T \otimes \dots$$

$$|\phi_2\rangle = \frac{1}{\sqrt{8}}(1,1,1,1,1,1,1,1)^T \otimes \dots$$

Example: 1st Iteration



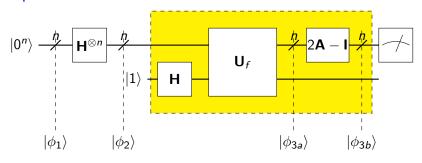
1st Iteration. Phase inversion:

$$|\phi_{3a}\rangle = \frac{1}{\sqrt{8}}(1,1,1,1,1,1,-1,1)^T \otimes \dots$$

Inversion about the mean $a = \frac{1}{\sqrt{8}} \frac{3}{4}$:

$$|\phi_{3b}\rangle = \frac{1}{\sqrt{8}}(\frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2})^T \otimes \dots$$

Example: 2nd Iteration



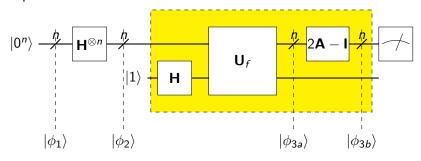
2nd Iteration. Phase inversion:

$$|\phi_{3a}\rangle = \frac{1}{\sqrt{8}}(\frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, -\frac{5}{2}, \frac{1}{2})^T \otimes \dots$$

Inversion about the mean $a = \frac{1}{\sqrt{8}} \frac{1}{8}$:

$$|\phi_{3b}\rangle = \frac{1}{\sqrt{8}}(-\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, \frac{11}{4}, -\frac{1}{4})^T \otimes \dots$$

Example: Measurement



Measuring the final state

$$|\phi_{3b}\rangle = \frac{1}{\sqrt{8}}(-\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, -\frac{1}{4}, \frac{11}{4}, -\frac{1}{4})^T \otimes \dots$$

in the standard base $|000\rangle$, $|001\rangle$,.... Will return (collapses to) one of the base vectors. We get with probability $|\frac{1}{1/8}\frac{11}{4}|^2=0.97227^2=0.94531$:

$$|110\rangle = (0,0,0,0,0,0,1,0)^T$$

Alternative Formulation: Oracle

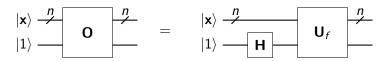
Some texts start immediately with a 'uniform' initial vector

$$|\psi\rangle = \frac{1}{\sqrt{n}} \sum_{x=0}^{n-1} |x\rangle$$

and use an oracle O for picking the element we look for.

$$\mathbf{0}: \sum_{\mathbf{x} \in \{0,1\}^n} \alpha_{\mathbf{x}} \, |\mathbf{x}\rangle \mapsto \sum_{\mathbf{x} \in \{0,1\}^n} (-1)^{f(\mathbf{x})} \alpha_{\mathbf{x}} \, |\mathbf{x}\rangle$$

In our presentation we obtained $|\psi\rangle$ explicitly by applying $\mathbf{H}^{\otimes n}$ while the oracle corresponds to the phase inversion:



Realisation of the Inversion about the Mean

As it is a unitary operator we know that $2\mathbf{A} - \mathbf{I}$ can be represented as a quantum gate. Concretely, some authors use the **conditional phase shift P**₀ to implement it.

$$\mathbf{P}_0: |\mathbf{x}
angle \mapsto \left\{ egin{array}{ll} |\mathbf{x}
angle & ext{if } \mathbf{x} = \mathbf{0}^n \\ -|\mathbf{x}
angle & ext{otherwise} \end{array}
ight.$$

As we can write $\mathbf{P}_0 = 2|0^n\rangle\langle 0^n| - \mathbf{I}$ we obtain

$$\mathbf{H}^{\otimes n}\mathbf{P}_{0}\mathbf{H}^{\otimes n}=\mathbf{H}^{\otimes n}(2\left|0^{n}\right\rangle\!\left\langle0^{n}\right|-\mathbf{I})\mathbf{H}^{\otimes n}=2\left|\psi\right\rangle\!\left\langle\psi\right|-\mathbf{I}=2\mathbf{A}-\mathbf{I}$$

The Grover operator $\mathbf{G} = ((2\mathbf{A} - \mathbf{I}) \otimes \mathbf{I})\mathbf{O}$ can be represented by

