

Identifying the Information Gain of a Quantum Measurement

Mario Berta, Joseph M. Renes, Mark M. Wilde - full version
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- Setup - making the question precise (information-theoretically)
- Main result - answer
- Proof ideas
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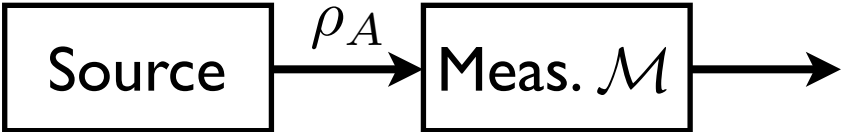
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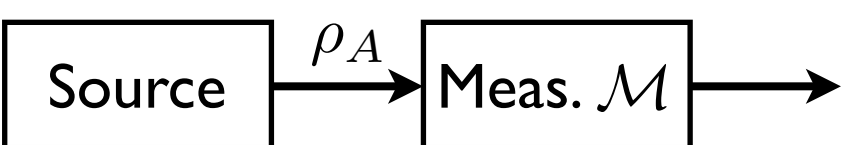
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$$\rho_X = \sum_x |x\rangle\langle x|_X \cdot \overbrace{\text{tr} \left[(M_A^x)^\dagger \underbrace{\rho_A}_{\rho_A^x} M_A^x \right]}^{p_x}$$

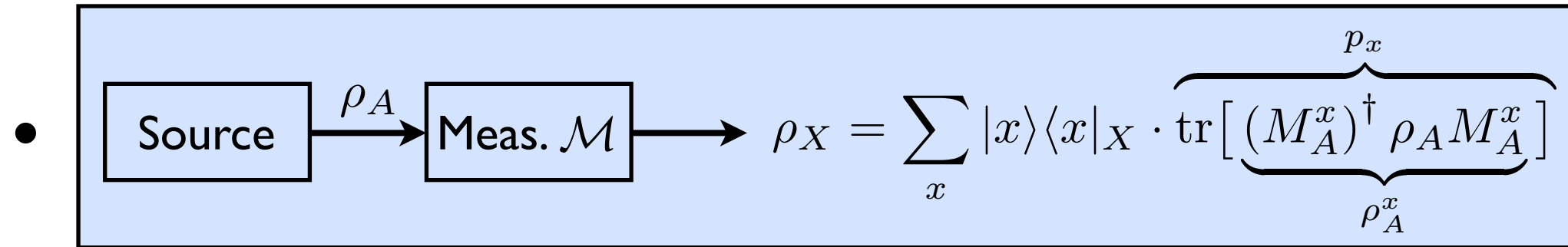
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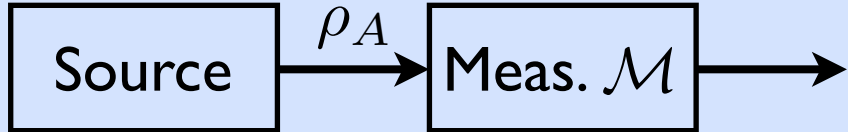
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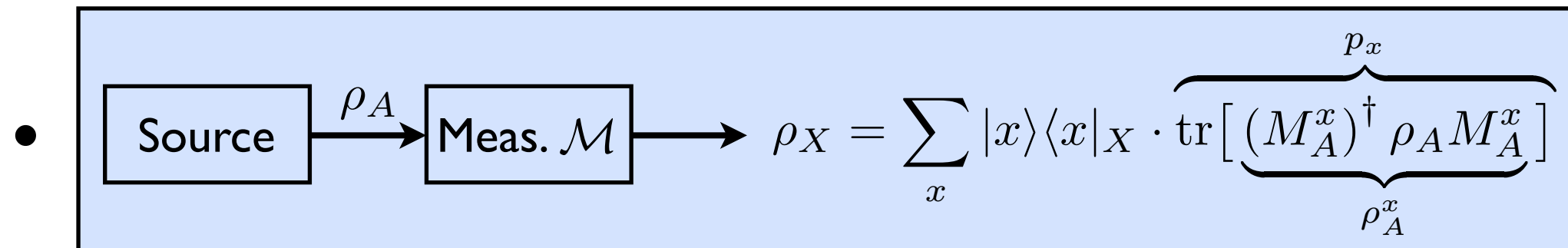
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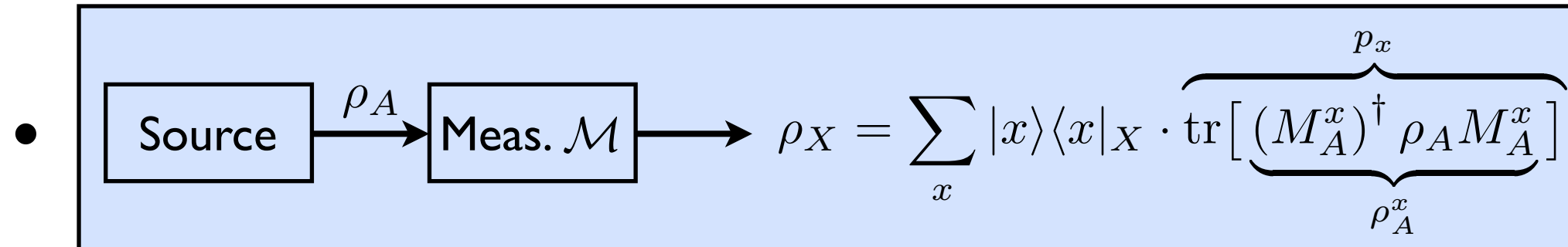
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$$I(\mathcal{M}) \stackrel{?}{=} H(X)_\rho \left(= - \sum_x p_x \log p_x \right)$$

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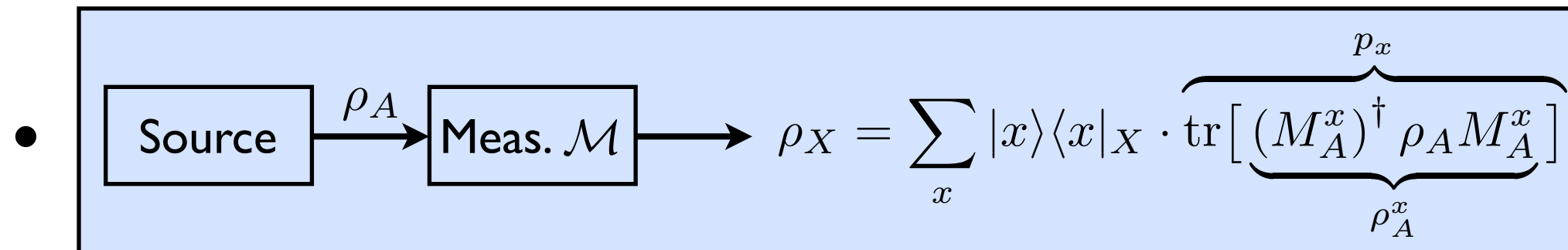
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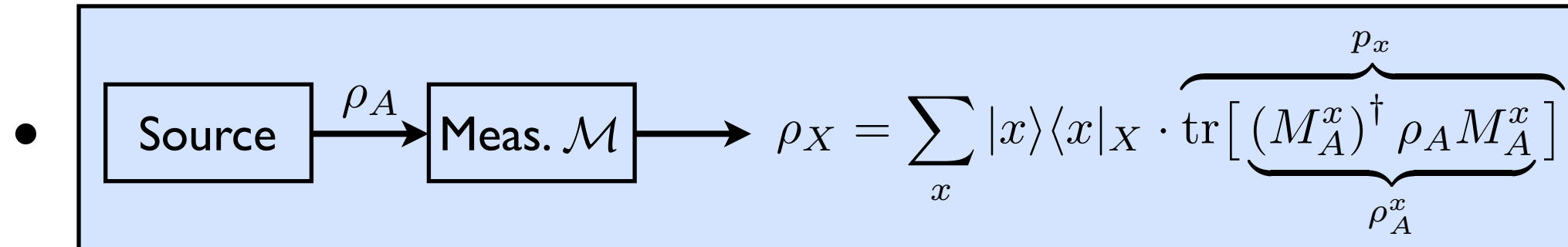
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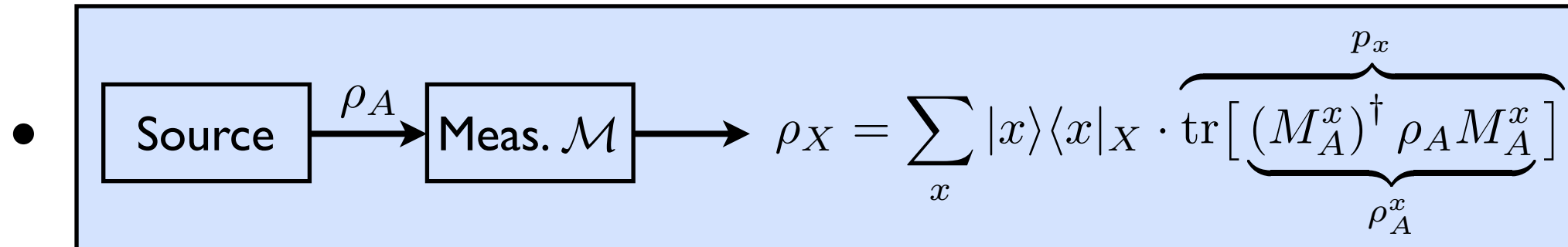
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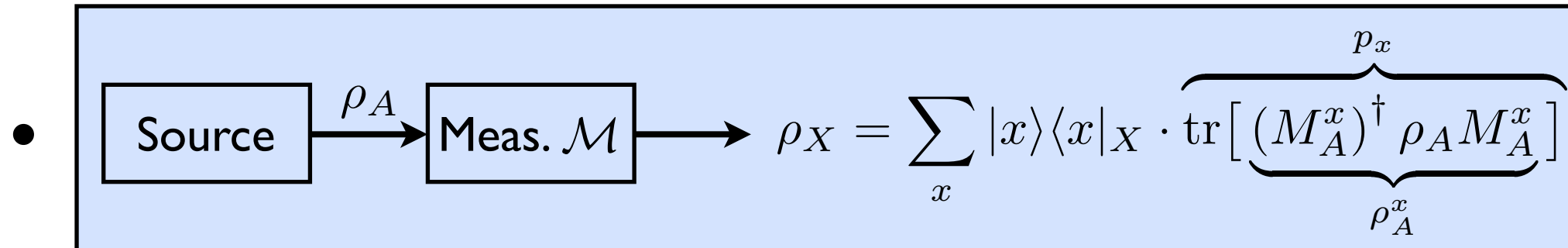
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- Information and disturbance in quantum measurements [3]:

$$I(\mathcal{M}) \stackrel{?}{=} I(X : R)_\omega (= H(X)_\omega + H(R)_\omega - H(XR)_\omega) \geq 0$$

$$\omega_{XR} = \sum_x |x\rangle\langle x|_X \otimes \text{tr}_A \left[(M_A^x \otimes 1_R)^\dagger \underbrace{\rho_{AR}}_{\text{pure}} (M_A^x \otimes 1_R) \right]$$

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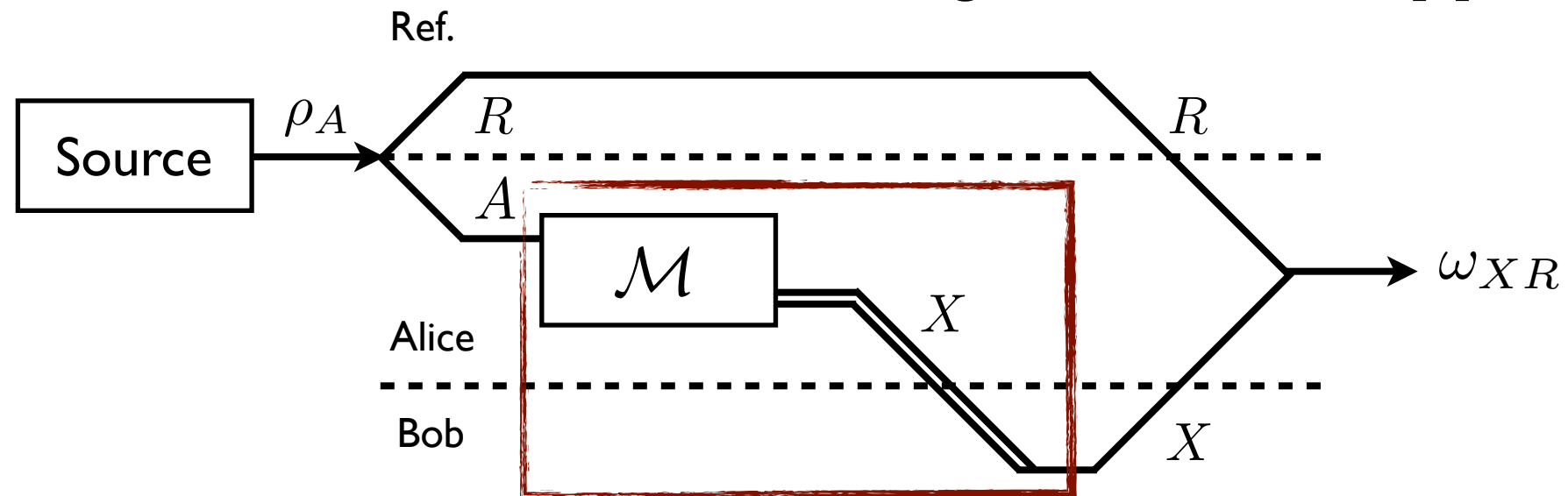
- Precise information-theoretic meaning in [4]!

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- Only output data that is correlated to the input should be counted as information gain! Information-theoretic idea of **simulating** the measurement [4]:

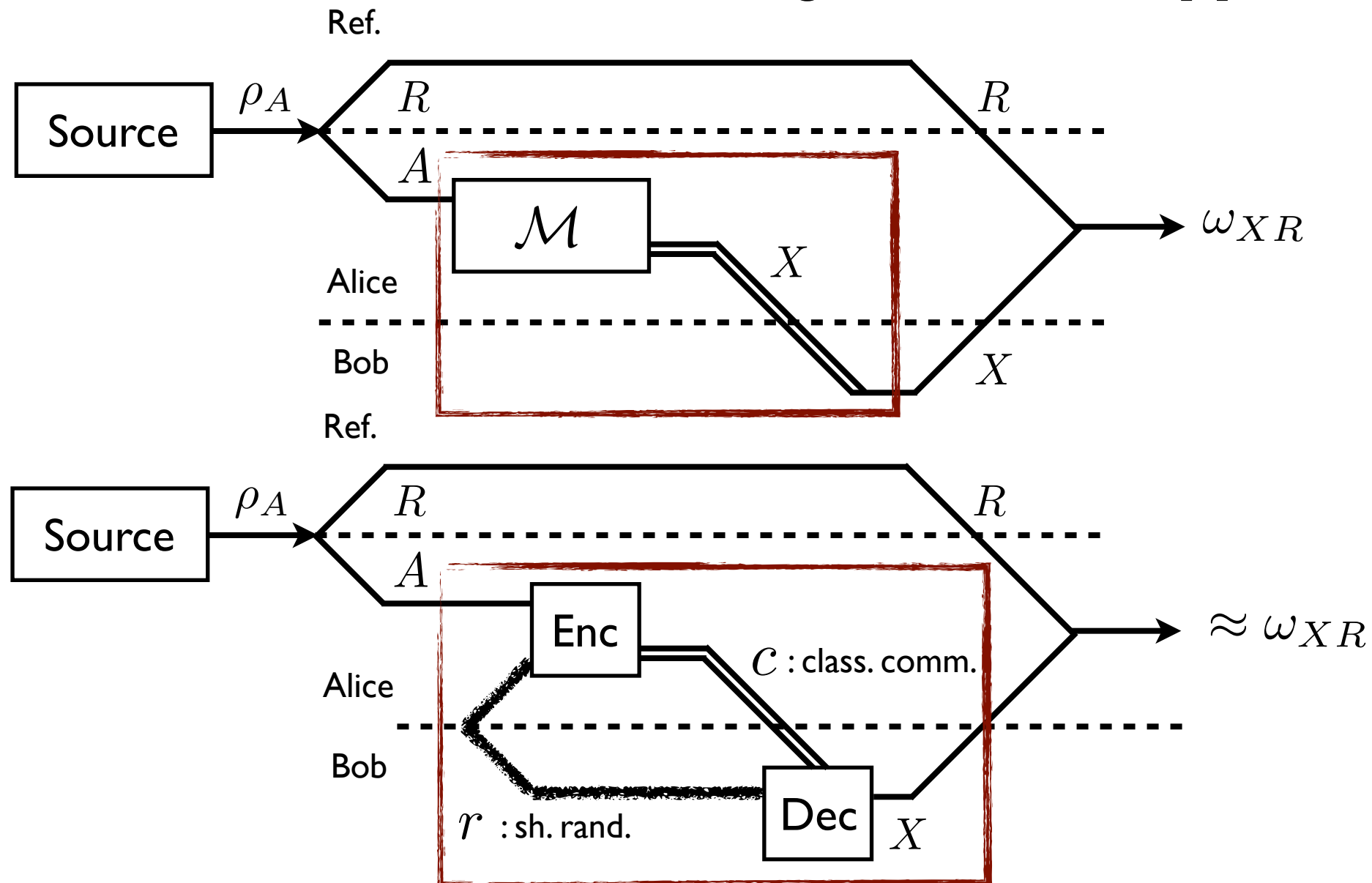
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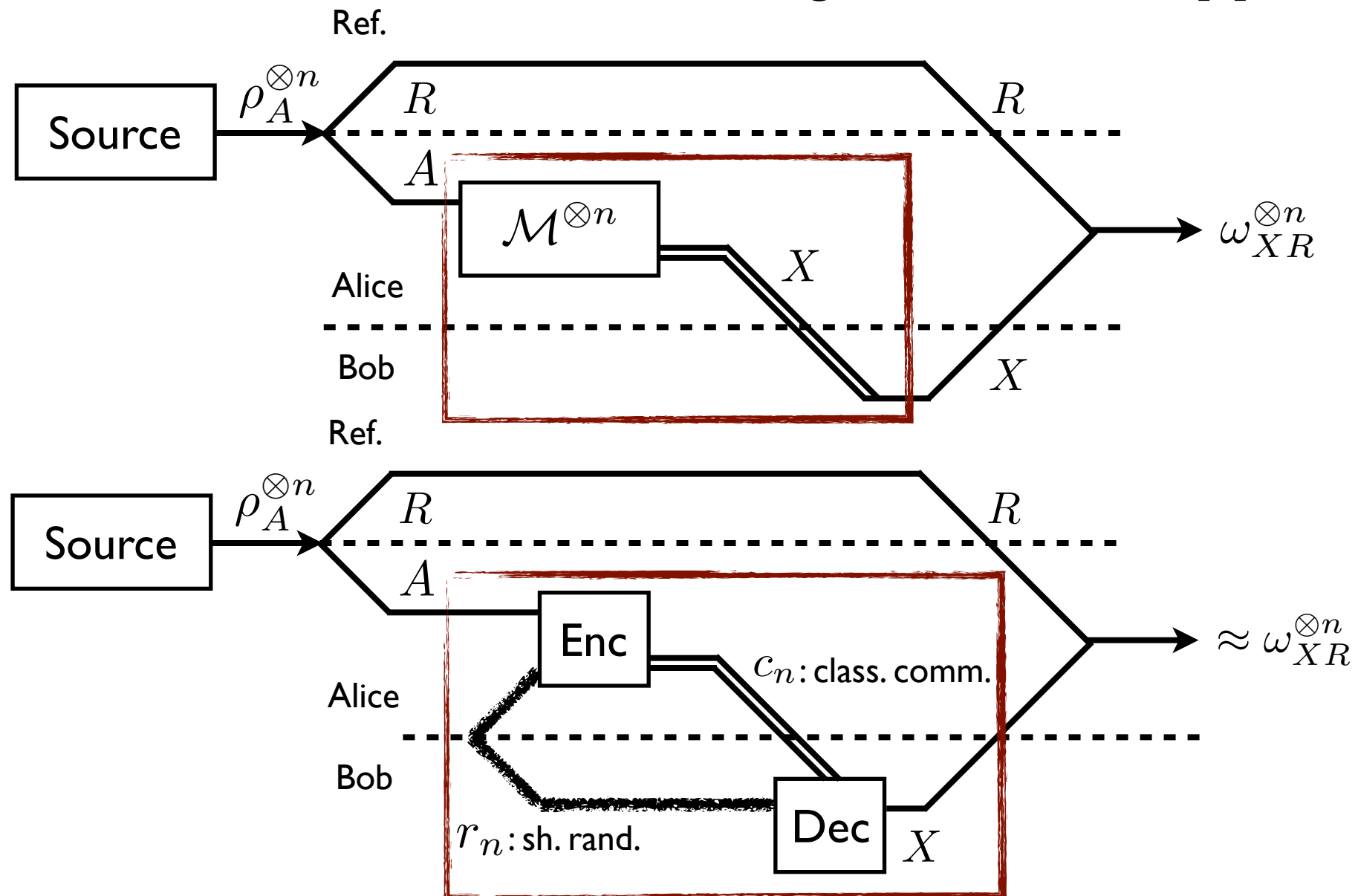
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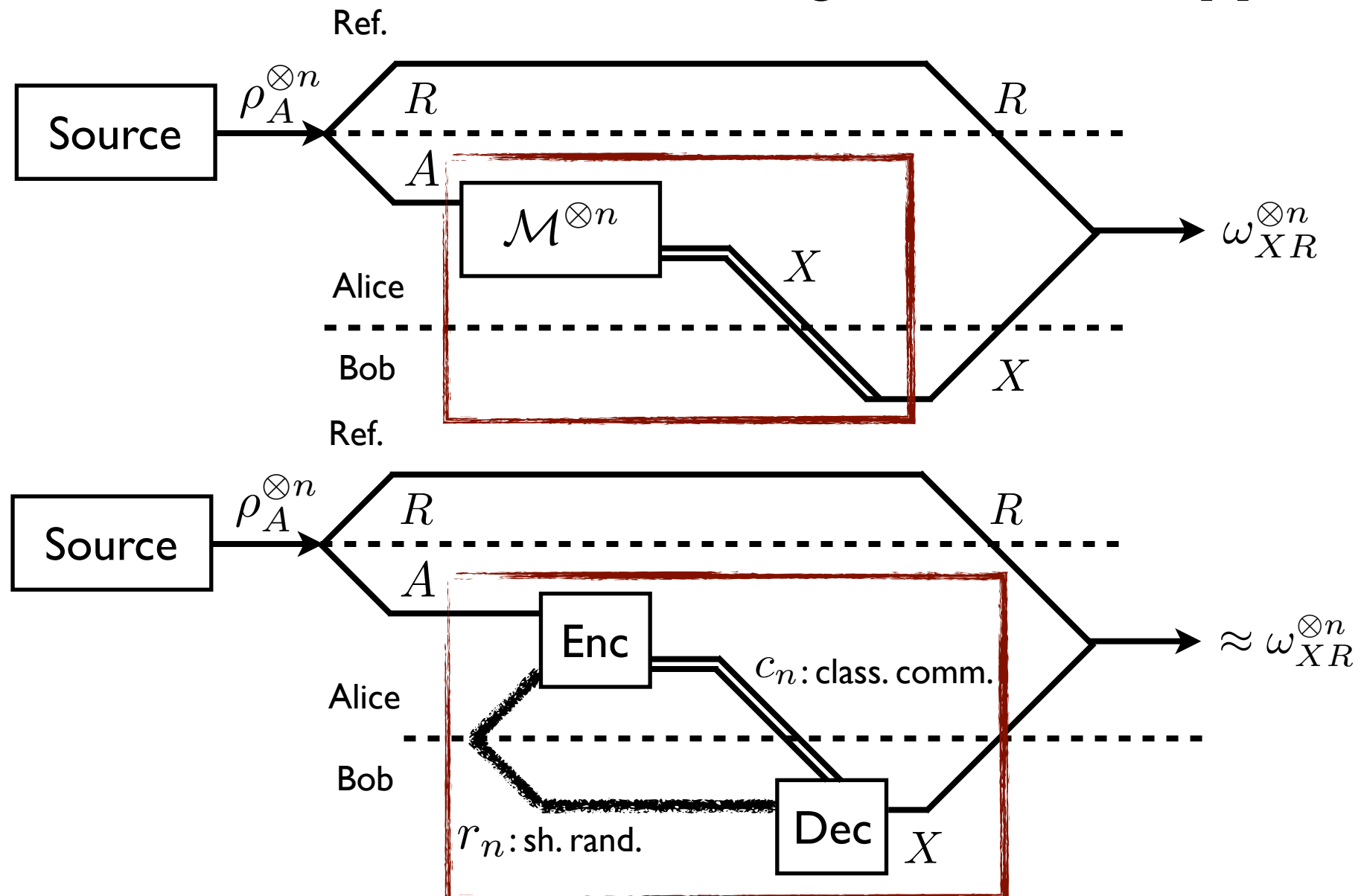
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- Main result in [4]:

$$c := \lim_{n \rightarrow \infty} \frac{c_n}{n} = I(X : R)_\omega$$

$$r := \lim_{n \rightarrow \infty} \frac{r_n}{n} = H(X|R)_\omega (= H(XR)_\omega - H(R)_\omega)$$

Setup (c)

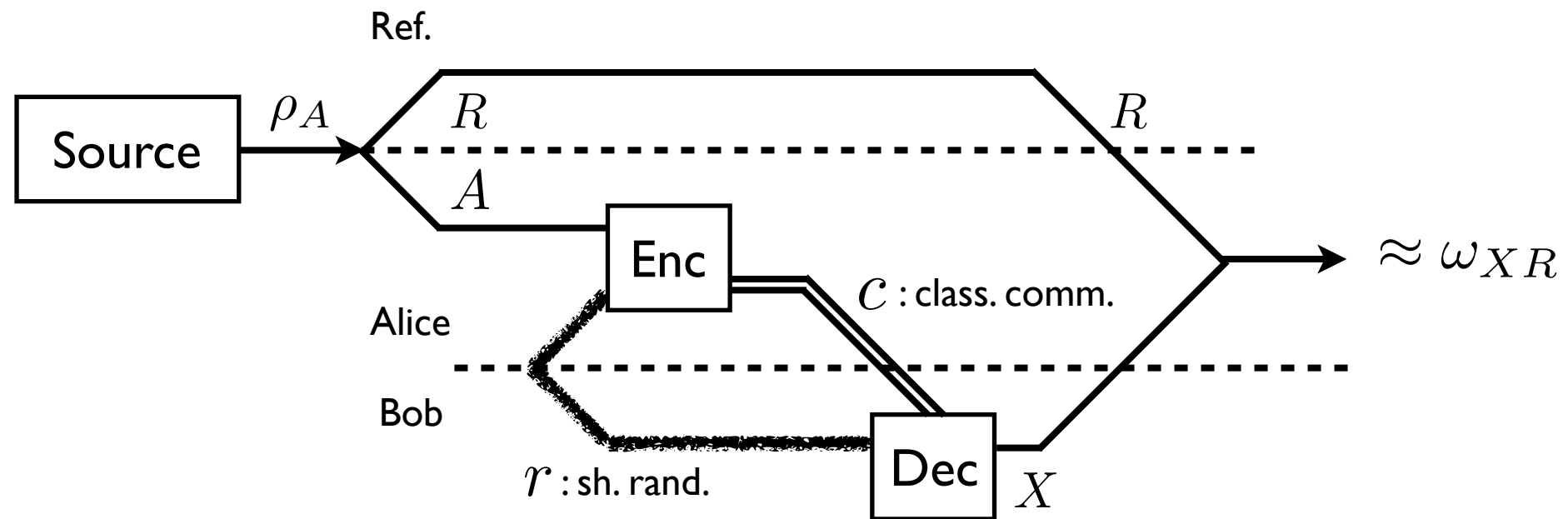
- An easy example for [4]:

$$\rho_A = \frac{1}{2}1_A \quad \mathcal{M} = \left\{ \frac{1}{2}|0\rangle\langle 0|, \frac{1}{2}|1\rangle\langle 1|, \frac{1}{2}|+\rangle\langle +|, \frac{1}{2}|-\rangle\langle -| \right\}$$

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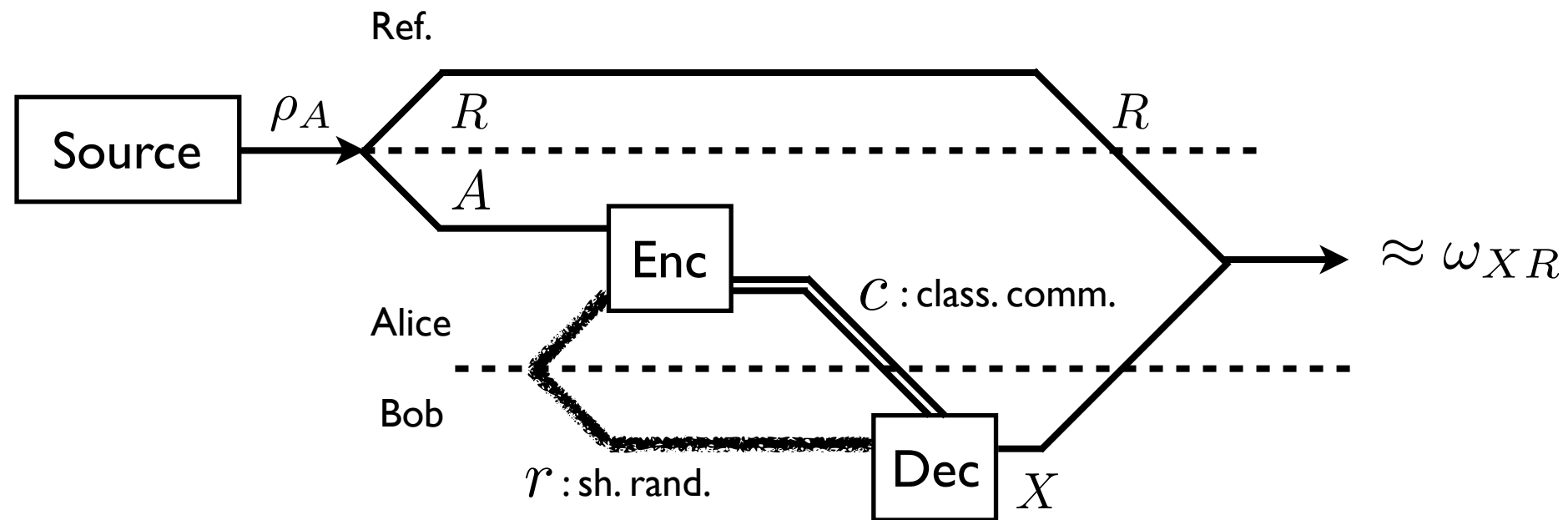
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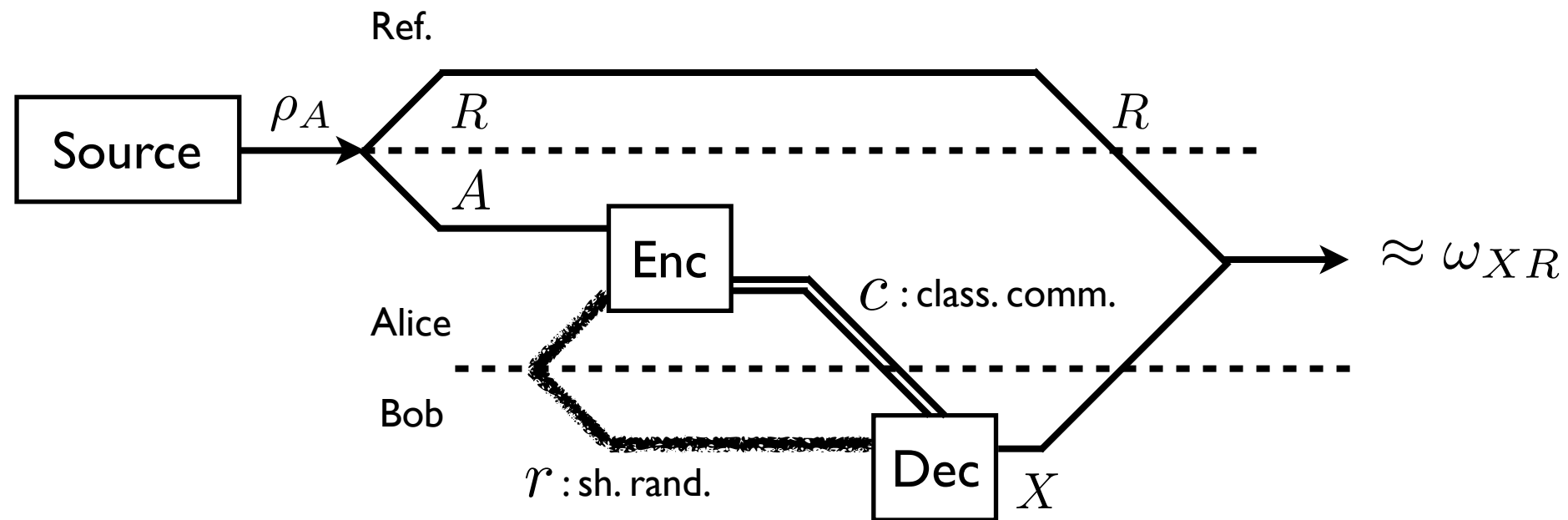


$$\rho_{AR} = |\rho\rangle\langle\rho|_{AR} \quad |\rho\rangle_{AR} = \frac{1}{\sqrt{2}} (|0\rangle_A \otimes |0\rangle_R + |1\rangle_A \otimes |1\rangle_R) = \frac{1}{\sqrt{2}} (|+\rangle_A \otimes |+\rangle_R + |-\rangle_A \otimes |-\rangle_R)$$

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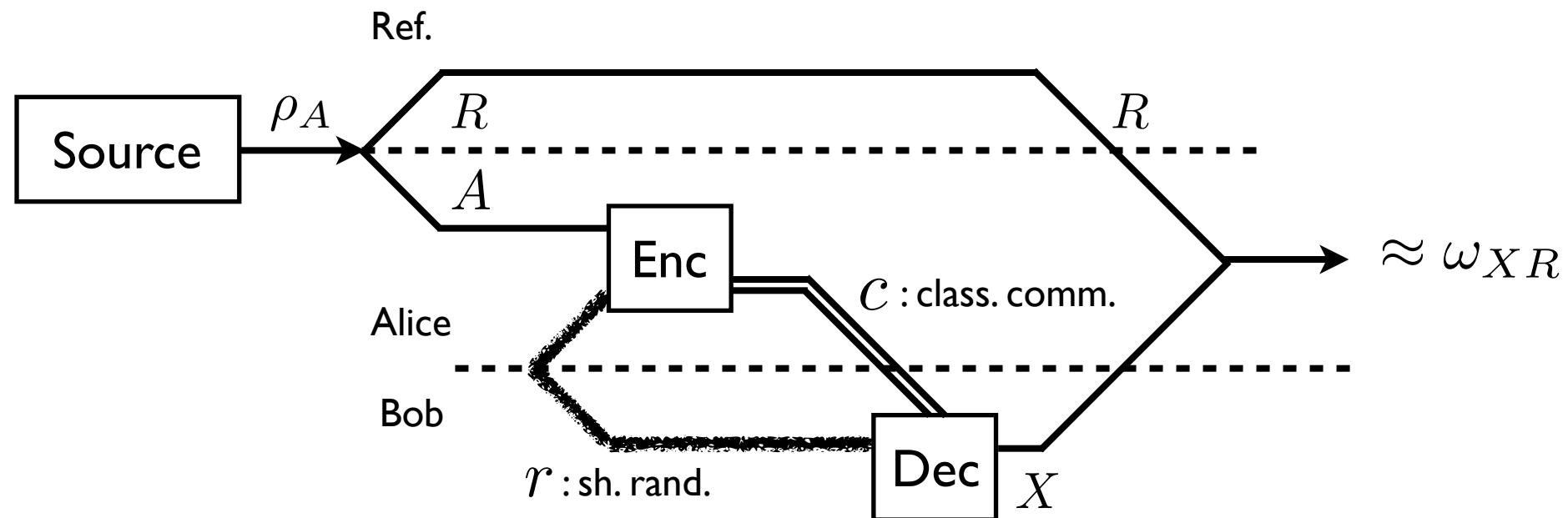
$$\omega_{XR} = \frac{1}{4} (|0\rangle\langle 0|_X \otimes |0\rangle\langle 0|_R + |1\rangle\langle 1|_X \otimes |1\rangle\langle 1|_R + |2\rangle\langle 2|_X \otimes |+\rangle\langle +|_R + |3\rangle\langle 3|_X \otimes |-\rangle\langle -|_R)$$

$$H(X)_\rho = 2 \quad I(X : R)_\omega = 1 \quad H(X|R)_\omega = 1$$

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- Apply according to shared randomness: $\sigma_X = \left\{ \frac{1}{2}|0\rangle\langle 0|, \frac{1}{2}|1\rangle\langle 1| \right\}$ $\sigma_Z = \left\{ \frac{1}{2}|+\rangle\langle +|, \frac{1}{2}|-\rangle\langle -| \right\}$

Question: *How much information (about the input) is gained by performing a given quantum measurement?*

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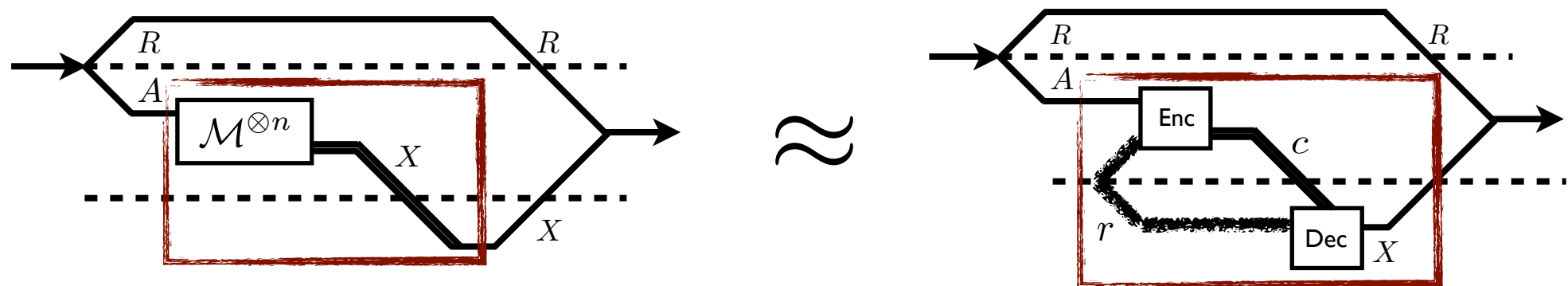
- Weaknesses of [4]: (i) only asymptotic iid (ii) **only for known fixed input** (protocol depends on input).
- We generalize (ii) to **arbitrarily varying sources** and even **entangled sources** (measurement simulation on general inputs):

$$\rho_A^{\otimes n} \text{ vs. } \rho_A^1 \otimes \rho_A^2 \otimes \cdots \otimes \rho_A^n \text{ vs. } \rho_{A^n}^n$$

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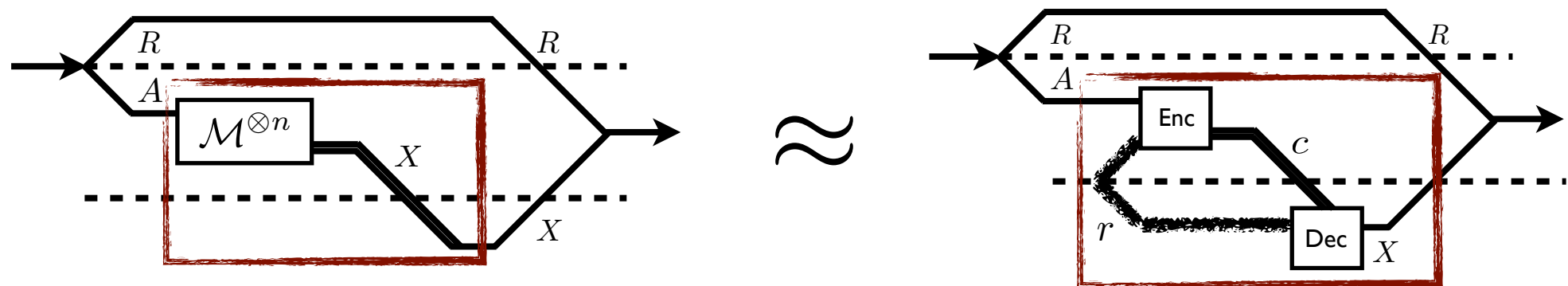
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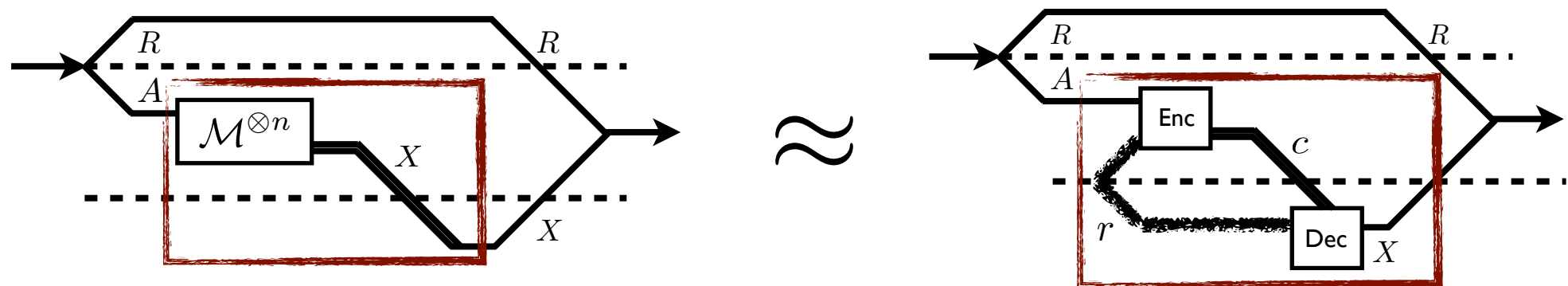


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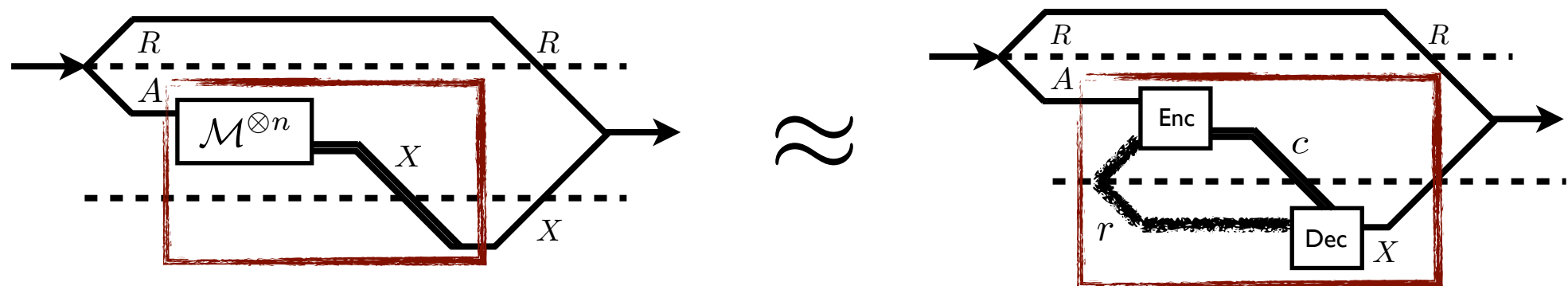
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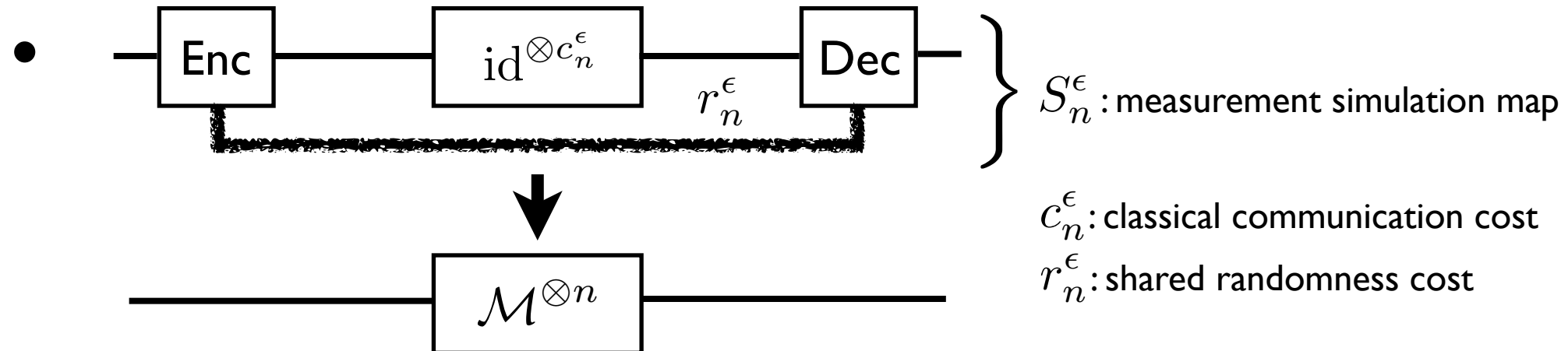
- By proving this, we also generalize (i) to the one-shot case: $c^{\epsilon} \approx I_{\text{max}}^{\epsilon}(X : R)_{\omega}$

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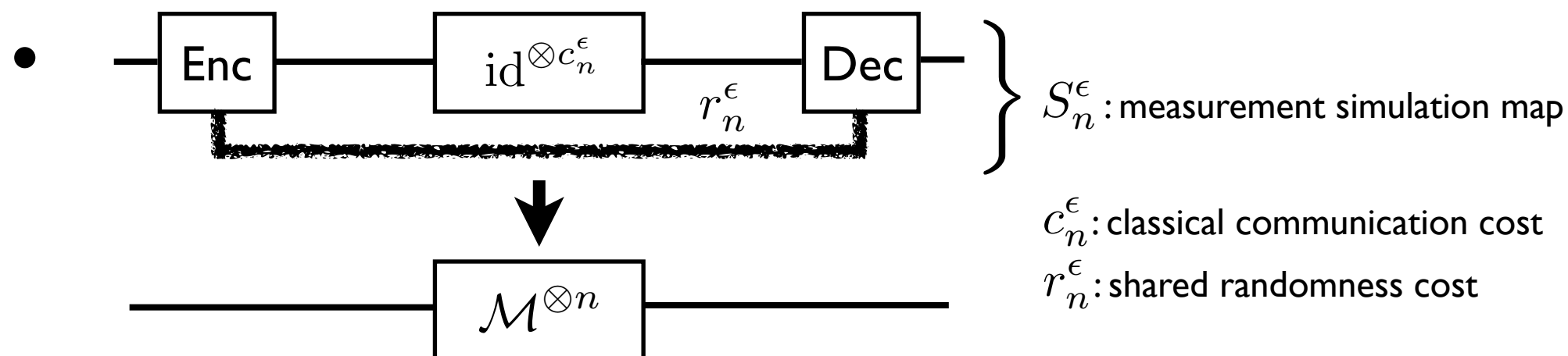
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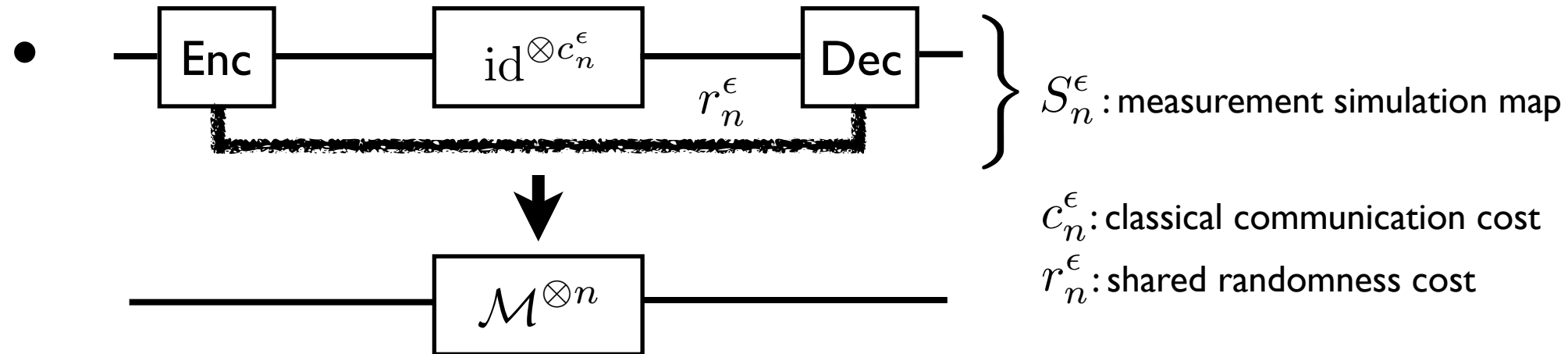


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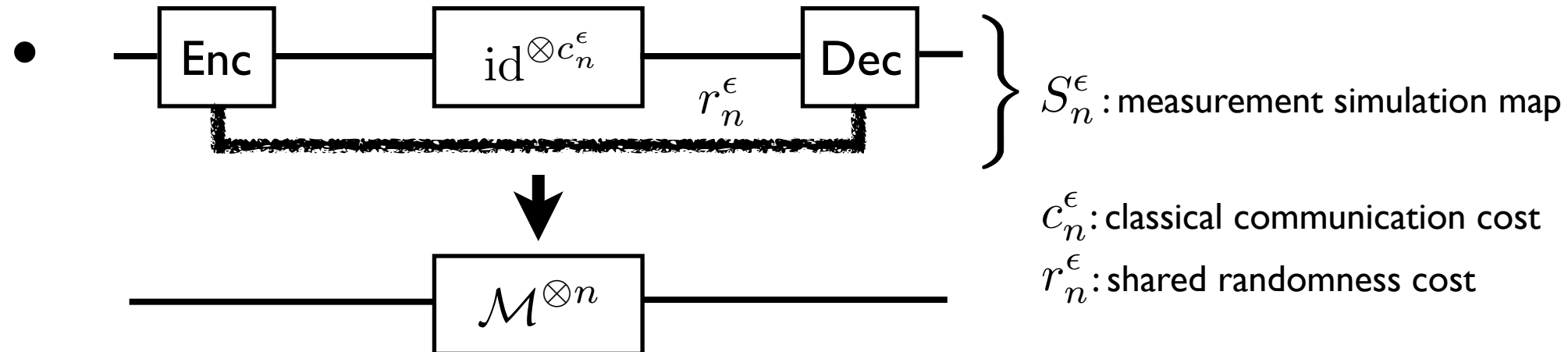
- In [4] for fixed iid input ρ_{AR} : $\left\| \left(\mathcal{M}_A^{\otimes n} \otimes \text{id}_R^{\otimes n} \right) \left(\rho_{AR}^{\otimes n} \right) - \left(S_n^\epsilon \otimes \text{id}_R^{\otimes n} \right) \left(\rho_{AR}^{\otimes n} \right) \right\|_1 \leq \epsilon$

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- Now: $\left\| \mathcal{M}_A^{\otimes n} - S_n^\epsilon \right\|_\diamond \leq \epsilon$ $\left\| \mathcal{M} \right\|_\diamond = \sup_{k \in \mathbb{N}} \sup_{\|\sigma\|_1 \leq 1} \left\| (\mathcal{M} \otimes \text{id}_k)(\sigma) \right\|_1$

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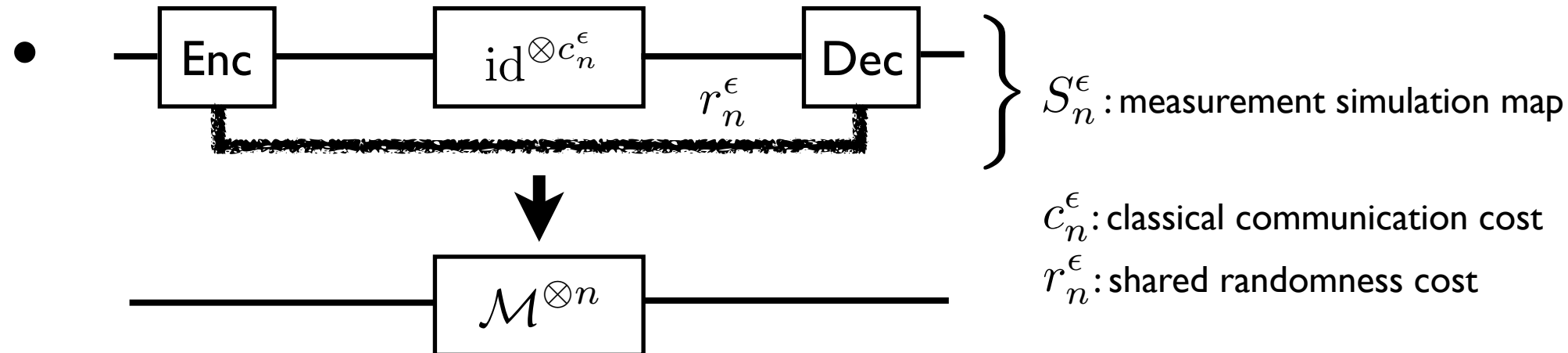
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- Post-selection technique for quantum channels [5]:**

$$\|\mathcal{M}_A^{\otimes n} - S_n^\epsilon\|_\diamond \leq \text{poly}(n) \cdot \|((\mathcal{M}_A^{\otimes n} - S_n^\epsilon) \otimes \text{id}_{R'}) (\zeta_{AR'}^n)\|_1$$

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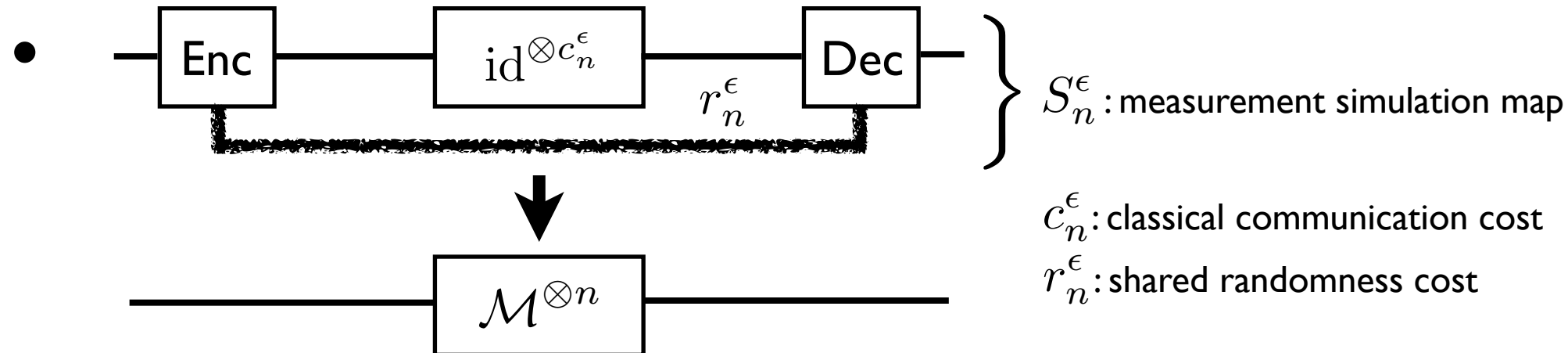
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$$\|\mathcal{M}_A^{\otimes n} - S_n^\epsilon\|_\diamond \leq \text{poly}(n) \cdot \|((\mathcal{M}_A^{\otimes n} - S_n^\epsilon) \otimes \text{id}_{R'}) (\zeta_{AR'}^n)\|_1$$

But: $\zeta_{AR'}^n$ purification of a special de Finetti state (a state which consists of n iid copies of a state on a single subsystem), i.e., no iid structure and [4] not applicable!

Proof Ideas



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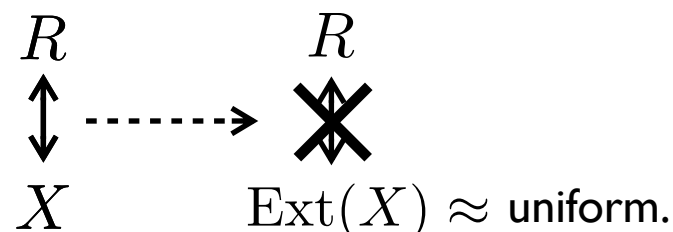
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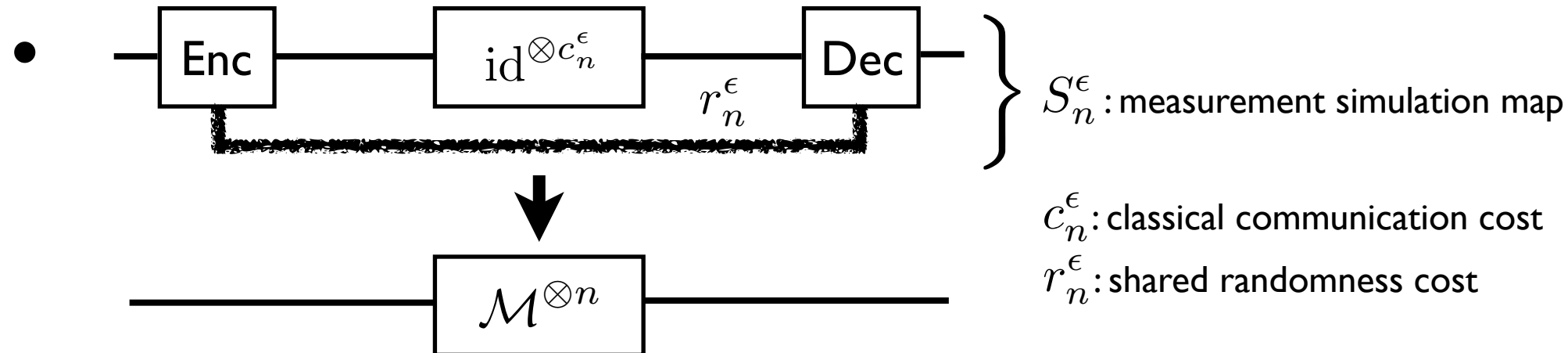
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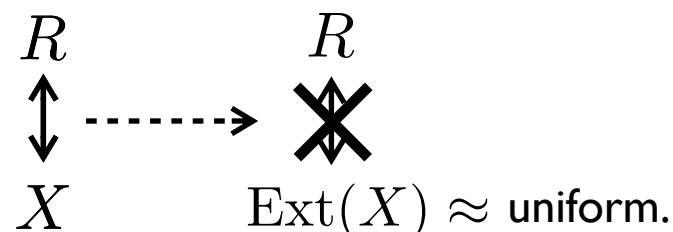
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Idea: extract all randomness from measurement data and only send the rest from Alice to Bob to simulate the measurement

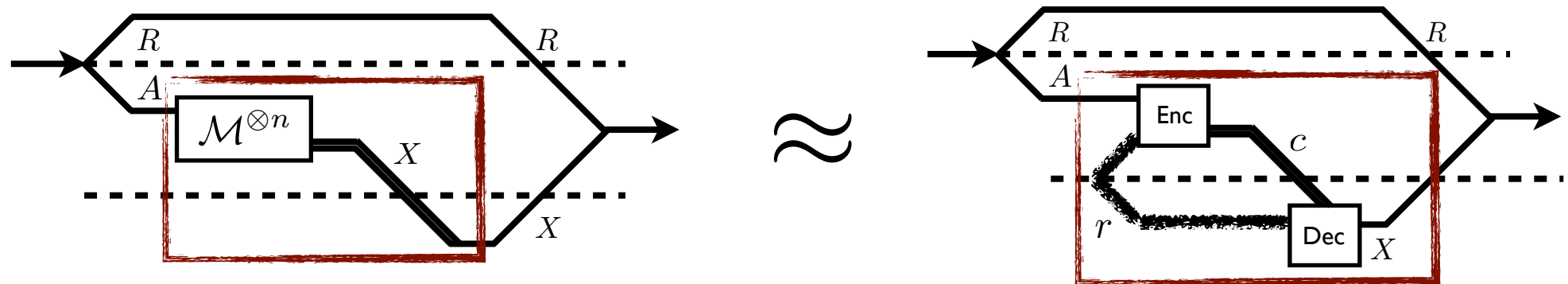
Question: How much information (about the input) is gained by performing a given quantum measurement?

Outline

- Setup - making the question precise (information-theoretically)
- Main result - answer
- Proof ideas
- **Conclusions**

Conclusions

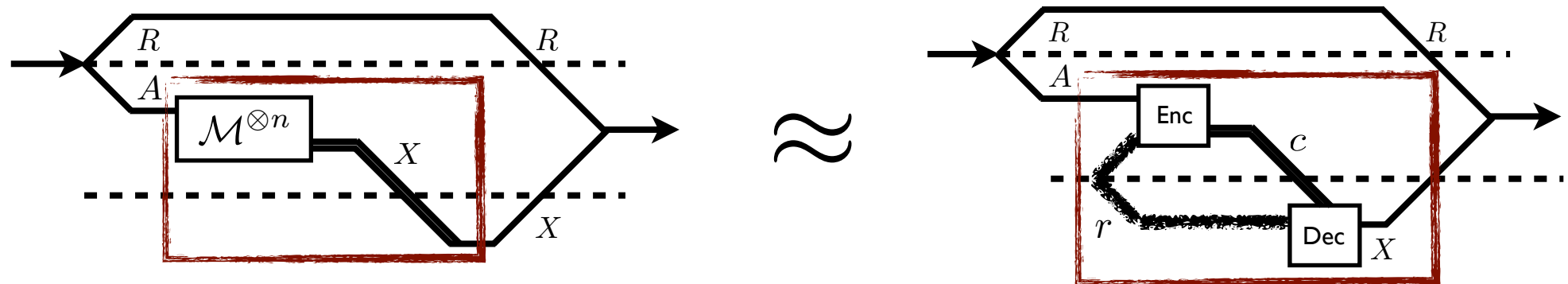
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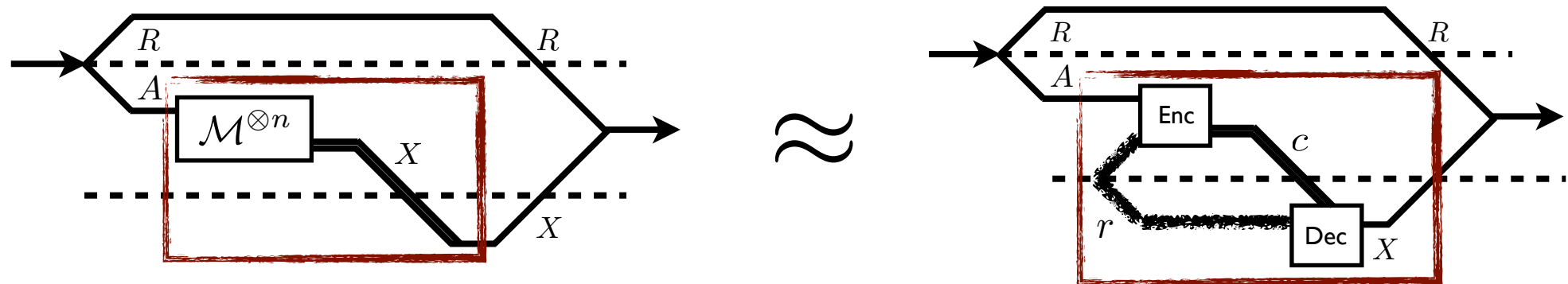
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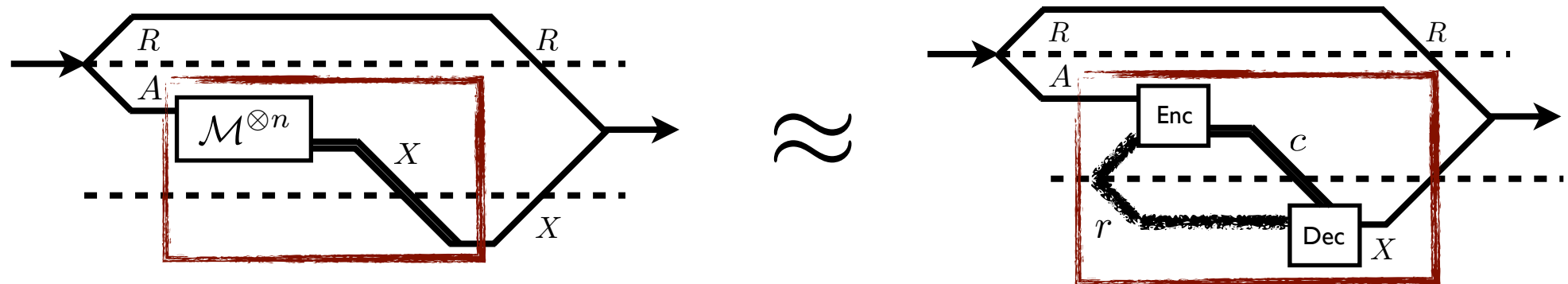
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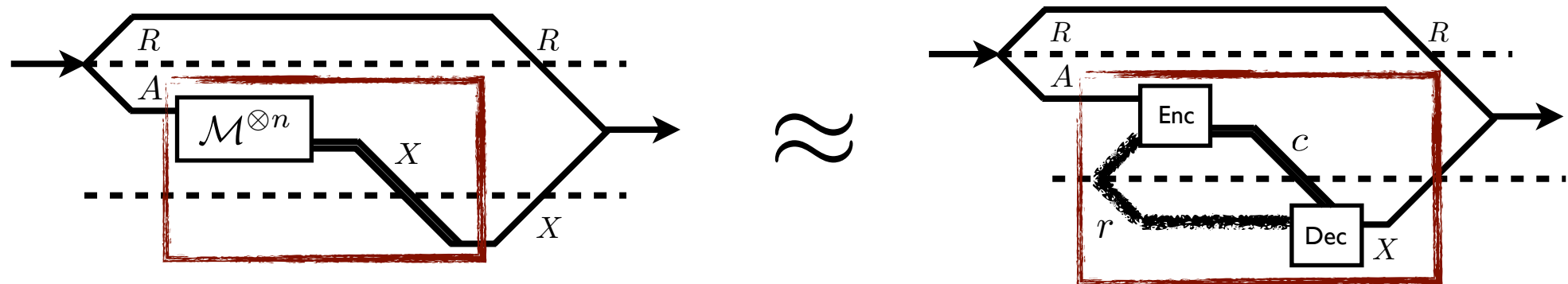
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- Extension: quantum instrument simulation
- Extension: explicit protocols

Thanks!